

THE SENTENTIAL CALCULUS USING RULE OF INFERENCE R_e

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Axiomatizations of the sentential calculus which use R_{mp} (*modus ponens*), have been shown equivalent to axiomatizations similar in all respects except that R_{mp} is replaced by the less restricted rule R_e (*rule of excision*)¹:

R_e . If S and $(\dots (S \supset S') \dots)$, then $(\dots S' \dots)$.

(Read: "If a formula 'S' and any formula 'T' containing '(S \supset S')' as a component are given, then one may proceed to a new formula 'T' got by substituting 'S' for '(S \supset S')' in 'T'."

We show here that R_e can be more powerful than R_{mp} , hence is independent.

I. The two following axioms, with R_e , yield the sentential calculus:

A1. $((s \supset (\phi \supset q)) \supset ((s \supset (q \supset r)) \supset (s \supset (\phi \supset r))))$

A2. $(r \supset ((-\neg q \supset -\phi) \supset (\phi \supset q)))$

For, from these, with R_e , we derive *1, *10, and *11 below; these constitute an axiom set which, with R_{mp} , yields the full sentential calculus²:

- *1. $((-\neg q \supset -\phi) \supset (\phi \supset q))$ [A2, A2 $r/(r \supset ((-\neg q \supset -\phi) \supset (\phi \supset q)))$, R_e]
- *2. $((\phi \supset q) \supset ((q \supset r) \supset (\phi \supset r)))$ [*1, A1 $s/((-\neg q \supset -\phi) \supset (\phi \supset q))$, R_e]
- *3. $(q \supset ((q \supset r) \supset r))$ [*1, *2 $\phi/((-\neg q \supset -\phi) \supset (\phi \supset q))$, R_e]
- *4. $(-\phi \supset -\phi)$ [*1, *3 $q/((-\neg q \supset -\phi) \supset (\phi \supset q))$, $r/-\phi$, R_e]
- *5. $((s \supset p) \supset ((s \supset (\phi \supset q)) \supset (s \supset q)))$ [*4, A1 $p/(-\phi \supset -\phi)$, q/p , r/q , R_e]
- *6. $(q \supset (\phi \supset p))$ [*4, A2 r/q , q/p , R_e]
- *7. $((((q \supset r) \supset r) \supset (\phi \supset r)) \supset (q \supset (\phi \supset r)))$ [*3, *2 p/q , $q/((q \supset r) \supset r)$, $r/(\phi \supset r)$, R_e]
- *8. $((\phi \supset (q \supset r)) \supset (((q \supset r) \supset r) \supset (\phi \supset r)))$ [*2 $q/((q \supset r) \supset r)$]
- *9. $((\phi \supset (q \supset r)) \supset (q \supset (\phi \supset r)))$ [*7, *8,
*2 $p/(\phi \supset (q \supset r))$, $q/(((q \supset r) \supset r) \supset (\phi \supset r))$, $r/(q \supset (\phi \supset r))$, R_e]
- *10. $(\phi \supset (q \supset p))$ [*6, *9 p/q , q/p , r/p , R_e]
- *11. $((s \supset (\phi \supset q)) \supset ((s \supset p) \supset (s \supset q)))$ [*5, *9 $p/(s \supset p)$, $q/(s \supset (\phi \supset q))$, $r/(s \supset q)$, R_e]

II. It is not possible to obtain ' $(\phi \supset p)$ ' from A1 and A2 if R_{mp} replaces R_e .

This is shown by assigning the accompanying interpretations, on which all provable theorems take only the value '0', to the primitive symbols ' \supset ' and ' $-$ '.

\supset	0	1	2	$-$	ϕ
0	0	2	2	2	0
1	0	2	2	1	1
2	0	0	0	0	2

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Received December 7, 1959.

¹ R. B. ANGELL, *On a less restricted type of rule of inference*, *Mind*, n.s. vol. LXIX (1960).

² J. ŁUKASIEWICZ and A. TARSKI, *Untersuchungen über den Aussagenkalkül*, *Comptes Rendus des Séances de la Société des Sciences et des Lettres de Varsovie*, Classe III, vol. 23 (1930), pp. 30–50.